Screening with damages and ordeals

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- ...and are a key factor in applicants' choice of development

- Thus, wait-times largely assume the role of prices:
 - They screen out low-value agents to balance supply and demand...
 - ...and "sort" participants into different types of units

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- However, wait-times have a curious screening property:
- They are more costly to households whose values for the units are higher
 - Each period of waiting deprives the household of the apartment's flow value
 - Thus, the cost of delaying receipt is multiplicative with value
- Other screening devices impose costs that are **separable** from values
 - E.g. differences in rent subsidies, application hassles...

Two kinds of screening devices

- I draw the distinction between damages and ordeals
- The cost of **damages** increases with the value for the good:
 - Waitlists and delays (through discounting or lost periods of use)
 - Damaged goods and usage restrictions (Deneckere and McAfee, 1996)
 - Network restrictions and changing claims rules in healthcare

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 - Waitlists and delays (through discounting or lost periods of use)
 - Damaged goods and usage restrictions (Deneckere and McAfee, 1996)
 - Network restrictions and changing claims rules in healthcare
- The cost of **ordeals** is separable from the value for the good:
 - Queues (Nichols et al., 1971)
 - Travelling to a distant office (Dupas et al., 2016)
 - Application hassles, bureaucracy (Deshpande and Li, 2019)

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- Result 1: with one good, never optimal to use damages
- Result 2: with two goods, using damages can be optimal
- Result 3: under regularity conditions, damages suboptimal even with two goods!

Also in the paper

- Heterogeneous costs of ordeals
- Monetary payments as *partially* wasteful screening
- Steady-state microfoundation for waitlist example
- Implications for affordable housing allocation

Model

Goods

- The designer distributes two kinds of goods, \boldsymbol{A} and \boldsymbol{B} ,
 - There is $\mu_A > 0$ of good A and $\mu_B > 0$ of good B

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- The designer distributes two kinds of goods, A and B,
 - There is $\mu_A > 0$ of good A and $\mu_B > 0$ of good B
- Agents' values for A and B are given by two-dimensional types (a,b)
 - Values (a,b) distributed according to F defined on $[0,1]^2$

Allocations

- The designer chooses a menu of damage and ordeal options for each of the goods
- That is, she chooses allocations of:
 - 1. Ordeals $t:[0,1]^2\to\mathbb{R}_+$
 - 2. Qualities $x:[0,1]^2 \to [0,1]$
 - 3. Goods: $y:[0,1]^2 \to \{A, B, \emptyset\}$
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- When x < 1, we say the good is **damaged**
- Type (a,b) who gets a good of quality x and completes an ordeal t gets utility:
 - $x \cdot a t$ if she gets A, $x \cdot b t$ if she gets B.

Designer's problem

- The designer maximizes total welfare:

$$W = \int U[a,b, (t,x,y)(a,b)] dF(a,b)$$

- She faces IC, IR and supply constraints:

for every
$$(a,b), (a',b') \in [0,1]^2, \quad U[a,b,(t,x,y)(a,b)] \ge U[a,b,(t,x,y)(a',b')]$$
 (IC)

for every
$$(a, b) \in [0, 1]^2$$
, $U[a, b, (t, x, y)(a, b)] \ge 0$

$$\int \mathbb{1}_{\mathrm{gets}\ A} \, \mathrm{d}F(a,b) \ \leq \ \mu_A, \qquad \int \mathbb{1}_{\mathrm{gets}\ B} \, \mathrm{d}F(a,b) \ \leq \ \mu_B$$

(S)

(IR)

One good case

One good case

- Suppose only good A is scarce, $\mu_A < 1$
- Good B is an unlimited outside option, $\mu_B = \infty$, with a common value b

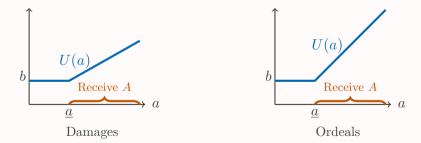
Proposition 1

Any mechanism that uses **damages**, so features x(a,b) < 1, is **Pareto dominated** by a mechanism that uses **only ordeals**.

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- ... so ordeals leave more rents to inframarginal takers of A!

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- ... but damages more costly to inframarginal types...
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- However, here, the A-good always goes to an upper interval of types
- With **2D** heterogeneity in values, there is **no fixed order!**
- Damages and ordeals **sort agents into goods** in different ways!

Two good case

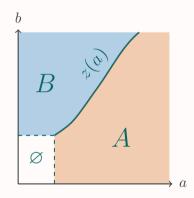
Two good case

- Consider the case where both goods are scarce: $\mu_A + \mu_B \leq 1$
- F, the distribution of values (a, b), has full support on $[0, 1]^2$

Feasible mechanisms

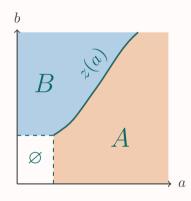
Who gets which good?

- When both goods come with ordeals, some types do not participate (\emptyset)
- The rest pick their favourite (ordeal, damage) option for one of the goods



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- When both goods come with ordeals, some types do not participate (\emptyset)
- The rest pick their favourite (ordeal, damage) option for one of the goods
- Types on the **boundary** z in different between their best options for both goods
- Types below z pick some option for A, types above z pick some option for B



Damages can be optimal

Ordeals and damages sort agents differently

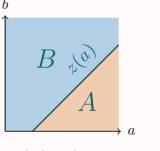
- Consider a mechanism with an **ordeal for each good**: c_A, c_B
- Then type (a, b) selects good A if $a c_A \ge b c_B$

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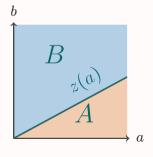
- Consider a mechanism with an **ordeal for each good**: c_A, c_B
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- Consider a mechanism which uses no ordeals but a damages good A to x = q
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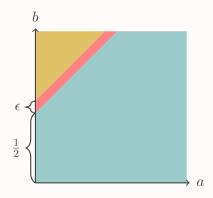
Only ordeals with $\mu_A + \mu_B = 1$



Only damages with $\mu_A + \mu_B = 1$

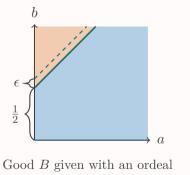
Example: damages can be optimal

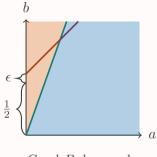
- Put mass ϵ on the **mustard** region...
- ... mass k on the **red** region...
- ...and mass $1-k-\epsilon$ on the **green** region
- Set supplies $\mu_A = 1 k \epsilon$, $\mu_B = k + \epsilon$



Example: damages can be optimal

- An "ordeal only" mechanism has $c_B=1/2,\ c_A=0$
- But this **eats away** almost all the surplus from getting B over A!





Good B damaged

- But a mechanism that damages B leaves surplus to agents close to the b-axis!

When are damages suboptimal?

When are damages suboptimal?

- 1. Consider piece-wise continuously differentiable $x:[0,1]^2 \to [0,1]$
- 2. The following are strictly increasing in one of a and b and non-decreasing in the other:

$$\frac{F_{A|B}(a|b)}{f_{A|B}(a|b)}, \quad \frac{F_{B|A}(b|a)}{f_{B|A}(b|a)},$$

Theorem 1

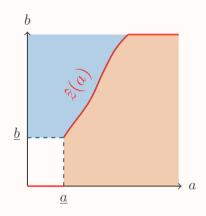
The optimal mechanism implements the efficient allocation of goods, and allocates both of them without damages. It posts a single ordeal for each good.

Proof strategy

Rewriting the objective

- Let $U_A: [0,1] \to \mathbb{R}_+$ be the indirect utility conditional on getting A
- Write total welfare as a function of U_A and the **extended boundary** \hat{z}

$$U_A(1) - \int_0^1 U_A'(a) \cdot F(a, \hat{z}(a)) da$$



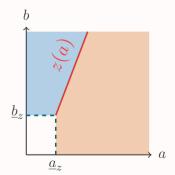
Extended boundary \hat{z} .

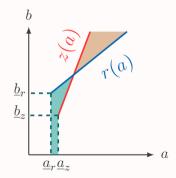
Proof strategy

$$U_A(1)-\int_0^1 U_A'(a)\cdot F(a,\hat{z}(a))da$$

- 1. Characterize implementable pairs (U_A, z)
- 2. Pick the optimal U_A for every fixed boundary z
- 3. Optimize over the space of optimally implemented boundaries z
- 4. Show the optimal boundary has a slope of 1 \rightarrow implementable without damages!

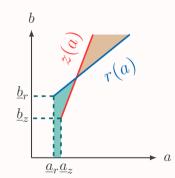
- Consider a linear boundary z with slope > 1...
- Our distributional assumptions will guarantee a less steep boundary is better
- Pick a less steep r such that z and r allocate the same amounts of A and B





- We can write the **difference in welfare** between r and z as:

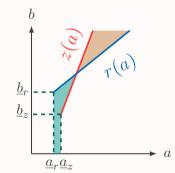
$$\Delta = (\underline{a}_z - \underline{a}_r) - \left(\int_{\overline{\mathcal{D}}} \frac{F_{A|B}(a \mid b)}{f_{A|B}(a \mid b)} f(a, b) d(a, b) - \int_{\underline{\mathcal{D}}} \frac{F_{A|B}(a \mid b)}{f_{A|B}(a \mid b)} f(a, b) d(a, b) \right).$$



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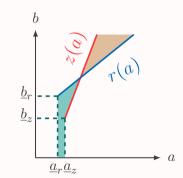
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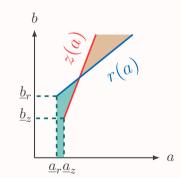
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- ...the masses in the **brown** and **green** regions are equal...



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- But $\underline{a}_z > \underline{a}_r \dots$
- ...the masses in the **brown** and **green** regions are equal...
- ... and $\frac{F_{A|B}(a|b)}{f_{A|B}(a|b)}$ is increasing in the \nearrow direction by assumption!



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- Screening devices differ in how they interact with agents' values
 - Damages impose costs that increase in one's value for the good
 - Ordeals impose costs that are separable from recipients' values
- Using damages is never optimal with only one kind of good
- And while they can be useful when many kinds of goods are offered...
- ...this is not the case for "regular" distributions

- Implications for **public housing** allocation?
 - Such programs often offer heterogeneous units, with different wait-times
 - Even if some wait-time is often inevitable in reality...
 - ... we should be worried about large **imbalances** in wait-times!
 - We should "sort" applicants using other instruments, e.g. by readjusting subsidies

Thank you!

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